

# Math 206A Lecture 11 Notes

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## 1 Polytopes and Permutahedra

### 1.1 Polytopes

**Definition 1.1.** A polytope  $P \subseteq \mathbb{R}^d$  is either of the following two equivalent things:

1.  $P = \text{conv}(X)$ ,  $|X| < \infty$ ,  $X \subseteq \mathbb{R}^d$ .
2.  $P = \bigcap H_i$  such that  $P$  is compact, where the  $H_i$  are half spaces.

**Definition 1.2.** The **dimension** of  $P$  is  $\dim(P) = \dim_{\mathbb{R}} \langle P \rangle$ , the affine subspace of  $\mathbb{R}^n$  spanned by  $P$ .

**Definition 1.3.** A **face**  $F \subseteq P$  is a subset of  $P$  such that there exists an affine subspace  $W \subseteq \mathbb{R}^d$  such that

1.  $F = P \cap W$ ,
2. there exists a half-space  $H$  such that  $P \subseteq H$ ,  $W \subseteq \partial H$ , and  $P \cap \partial H = F$ .<sup>1</sup>

**Example 1.1.** Let  $C_3$  be the cube in  $\mathbb{R}^3$ . Then  $C_3 = \text{conv}(\{(\pm 1, \pm 1, \pm 1)\})$ . On the other hand,  $C_3 = \bigcap_{i=1}^3 \{x : x_i \leq 1\} \cap \bigcap_{i=1}^3 \{x : x_i \geq -1\}$ . The faces are  $C_3 \cap \{x : x_i = \pm 1\}$ .

**Definition 1.4.** An **edge** is a 1-dimensional face. A **vertex** is a 0-dimensional face. A **facet** is a  $(d - 1)$ -dimensional face.

**Definition 1.5.** The **graph**  $\Gamma(P)$  of a polytope is a graph  $\Gamma = (V, E)$ , where  $V(P)$  is the set of vertices of  $P$  and  $E(P)$  is the set of edges of  $P$ .

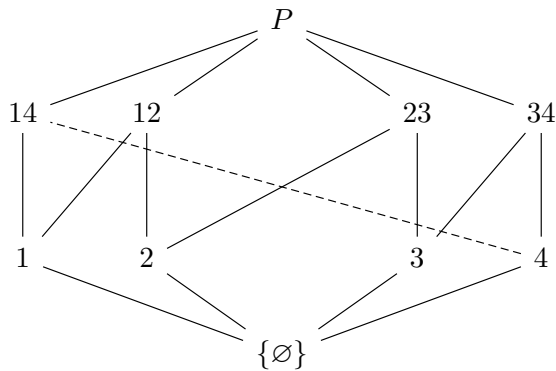
**Definition 1.6.** The **face lattice**  $\alpha(P)$  is the partially ordered set of faces of  $P$ , ordered by inclusion.

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<sup>1</sup>This second condition implies the first, so you should really think of it as a clarification of the previous condition.

This is a lattice because the meet of  $F$  and  $F'$  is  $F \cap F'$ , and the join of  $F$  and  $F'$  is  $\langle F \cup F' \rangle \cap P$ .

**Example 1.2.** Let  $P$  be a square in  $\mathbb{R}^2$ . Then  $\Gamma(P)$  is the graph of the boundary of the square, and  $\alpha(P)$  is



Here is a theorem we will prove later.

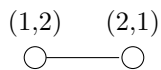
**Theorem 1.1** (Blind-Mani). *If  $P \subseteq \mathbb{R}^d$  is simple, then  $\Gamma$ , the graph of  $P$  determines the face lattice of  $P$ .*

## 1.2 Permutahedra

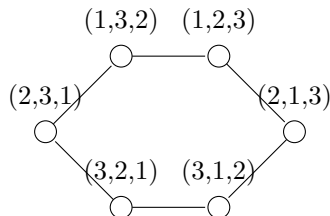
**Definition 1.7.** This **permutahedron** is  $P = \text{conv}(\{(\sigma(1), \dots, \sigma(n)) : \sigma \in S_n\})$ .

Observe that  $\dim(P) = n - 1$ .

**Example 1.3.** For  $n = 2$ , the permutahedron is



For  $n = 3$ , we have



**Proposition 1.1.**  $\Gamma(P_n) \cong \text{Cay}(S_n, R_n)$ , where  $R_n$  is the set of transpositions  $(i j)$  with  $1 \leq i, j \leq n$  with a left action.

Observe that  $P_n$  is simple. In particular, we can figure out the  $\mathcal{F}$ -vector. Consider a linear functional  $\varphi : \mathbb{R}^n \rightarrow \mathbb{R}$  nonconstant on edges where

$$\varphi(x_1, \dots, x_n) = x_1 + \varepsilon x_2 + \varepsilon^2 x_3 + \dots + \varepsilon^{n-1} x_n.$$

Then  $\text{ind}(\sigma)$  is the number of  $i \in \{1, \dots, n-1\}$  such that  $\varphi((i \ i+1)\sigma) > \varphi(\sigma)$ . This is the number of  $i$  such that  $\sigma^{-1}(i) < \sigma^{-1}(i+1)$ . So  $g_k = h_k^{(\varphi)}$  is the number of  $\sigma \in S_N$  such that  $\sigma^{-1}$  has  $k$  ascents.

Let  $A(n, K)$  be the number of  $\sigma \in S_n$  with exactly  $k$  ascents.<sup>2</sup> We can prove the following proposition.

**Proposition 1.2.**

$$A(n, k) = (n - K)A(n - 1, k - 1) + (k + 1)A(n - 1, k)$$

**Example 1.4.** This is called the **Birkhoff polytope**. It is the set of matrices of non-negative entries such that the sum of the rows and columns are all 1. Formally, this is  $B_n = \bigcap_{i=1}^n \bigcap_{j=1}^n \{x : x_{i,j} \geq 0\} \cap \bigcap_{j=1}^n \{x : \sum_{i=1}^n x_{i,j} = 1\} \cap \bigcap_{i=1}^n \{x : \sum_{j=1}^n x_{i,j} = 1\}$ .

**Theorem 1.2.**  $V(B_n) = \{\text{Mat}(\sigma) : \sigma \in S_n\}$ .  $E(B_n) = \{(\sigma, w\sigma) : w \in S_n \text{ is a cycle}\}$

**Corollary 1.1.**  $\deg_{\Gamma}(v)$  is the number of cycles in  $S_n$ , and  $\dim(B_n) = (n - 1)^2$ .

Question: Is  $f_i(B)$  computable in polynomial time?

**Theorem 1.3** (Pak). *Let  $Q_n$  be the set of such matrices but with dimension  $n \times (n + 1)$ . Then  $f_i(Q_n)$  can be computed in polynomial time.*

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<sup>2</sup>The bivariate generating function for  $A(n, k)$  has a nice form.